# Algorithm Design: Fairness Versus Accuracy

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# background

algorithms are used to guide many high-stakes decisions

- who should receive a medical treatment?
- who should receive a loan?
- who should receive bail?
- who should receive employment?

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recent empirical evidence that these algorithms often have errors that vary systematically across subgroups of the population

- patients assigned to same risk score have substantially different actual health risks depending on race (Obermeyer et al., 2019)
- false positive rate of algorithm used to predict criminal reoffense twice as high for Black defendants (Angwin and Larson, 2016)

#### fairness vs. accuracy

- algorithms increasingly optimized not only for accuracy but also "fairness" (equalizing error rates across groups)
- what is the tradeoff between fairness and accuracy?
- we introduce a "fairness-accuracy frontier" that ranges across a broad class of preferences/optimization criteria
- results characterize how this frontier depends on statistical properties of the inputs to the algorithm
  - whether the inputs reveal the group identity
  - whether the inputs are group-balanced
- in paper but will skip in talk: "input design" problem

# framework

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- each subject is described by three variables:
  - **type** Y taking values in  $\mathcal{Y}$  (e.g. need for medical procedure)
  - group  $G \in \mathcal{G} = \{r, b\}$  (e.g. race)
  - covariate vector X taking values in X
     (e.g. image scans, # past hospital visits, blood tests)

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- X is observed by the designer, Y and G are not directly observed (but may be revealed by X)
- $\bullet$  in the population,  $(\mathit{Y}, \mathit{G}, \mathit{X}) \sim \mathbb{P}$

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the designer chooses an algorithm

 $a:\mathcal{X}\to\Delta(\mathcal{D})$ 

for determining (distributions over) decisions based on the observed covariate vector

#### group errors

fix a loss function  $\ell:\mathcal{D}\times\mathcal{Y}\times\mathcal{G}\rightarrow\mathbb{R}$ 

 $\bullet$  can interpret  $\ell$  as measure of inaccuracy or as the disutility of a given subject

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#### Definition

the **error** for group  $g \in \mathcal{G}$  given algorithm *a* is

$$e_{g}\left(a
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i.e., the average/expected loss for subjects in group g

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- improving accuracy: lowering er and eb
- improving fairness: lowering  $|e_r e_b|$

### special cases

this approach nests several existing fairness metrics:

**example:** equality of false positive rates corresponds to  $e_r(a) = e_b(a)$  with

$$\ell(d, y) = \begin{cases} 1 & \text{if } (d, y) = (1, 0) \\ 0 & \text{otherwise} \end{cases}$$

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example: algorithm a satisfies equalized odds if

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this corresponds to  $e_r(a) = e_b(a)$  with

$$\ell(d, y, g) = \begin{cases} \frac{P(Y = y)}{P(Y = y \mid G = g)} & \text{if } d = 1\\ 0 & \text{otherwise} \end{cases}$$

## preferences

#### Definition (fairness-accuracy (FA) dominance)

let  $>_{F\!A}$  be the partial order on  $\mathbb{R}^2$  satisfying  $(e_r,e_b)>_{F\!A}(e_r',e_b')$  if

$$\underbrace{e_r \leq e_r', \quad e_b \leq e_b'}_{\text{higher accuracy}} \quad \text{and} \quad \underbrace{|e_r - e_b| \leq |e_r' - e_b'|}_{\text{higher fairness}}$$

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#### Definition

a fairness-accuracy preference  $\succeq$  is any total order on  $\mathbb{R}^2$  such that  $e \succ e'$  whenever  $e >_{F\!A} e'$ 

• utilitarian/bayes-optimal:

 $w_u(e_r, e_b) = -p_r e_r - p_b e_b$ 

where  $p_r$  and  $p_b$  are the proportions of either group

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- **3** rawlsian/group DRO: first order errors by  $-\max\{e_r, e_b\}$ , and then break ties using  $w_u$

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**onstrained optimization** (e.g., Hardt et al., 2016):

 $\min_{a\in\mathcal{A}_X} \quad p_re_r(a)+p_be_b(a) \quad ext{ s.t. } |e_r(a)-e_b(a)|\leq arepsilon$ 

# fairness-accuracy frontier

#### Definition

#### the **feasible set** given X is

$$\mathcal{E}(X) := \{(e_r(a), e_b(a)) : a \in \mathcal{A}_X\}$$

where  $\mathcal{A}_X$  is the set of all algorithms  $a : \mathcal{X} \to \Delta(\mathcal{D})$ 

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 $\mathcal{F}(X) := \left\{ e \in \mathcal{E}(X) : \nexists \ e' \in \mathcal{E}(X) \ \text{s.t.} \ e' >_{FA} e \right\}$ 

• includes optimal points across all fairness-accuracy preferences

# characterizing the

# fairness-accuracy frontier

## feasible set of group error pairs

**lemma**: for any X, the feasible set  $\mathcal{E}(X)$  is compact and convex (if  $\mathcal{X}$  is finite, it is a convex polygon)



group-optimal points:

$$R_X := \arg\min_{e \in \mathcal{E}(X)} e_r \qquad B_X := \arg\min_{e \in \mathcal{E}(X)} e_b$$

fairness-maximizing point:

$$F_X := \arg\min_{e \in \mathcal{E}(X)} |e_r - e_b|$$

(break all ties in favor of aggregate accuracy)











### group-skewed vs group-balanced

#### Definition

#### covariate vector X is

• *r*-skewed if  $e_r < e_b$  at  $R_X$  and  $e_r \leq e_b$  at  $B_X$ 

"group r's error is lower both at group r's favorite point and also at group b's favorite point"

- *b*-skewed if  $e_b < e_r$  at  $B_X$  and  $e_b \le e_r$  at  $R_X$
- group-balanced otherwise

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group-skew can emerge in practice (for example) if the inputs in X are systematically more informative about Y for one group

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#### Theorem

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- $R_X$  and  $B_X$  if X is group-balanced
- $G_X$  and  $F_X$  if X is g-skewed (usual Pareto frontier + more)



## strong fairness-accuracy conflict

**corollary:** there exist  $e, e' \in \mathcal{F}(X)$  satisfying

$$e_r \leq e_r', \quad e_b \leq e_b', \quad |e_r-e_b| > |e_r'-e_b'|$$

e.g., 
$$e = (1/3, 1/4)$$
,  $e' = (1/2, 1/2)$ 

if and only if X is group-skewed



"I have a policy proposal, which would decrease accuracy for both groups, but increase fairness."

ACADEMIC

"Are the inputs to your algorithm group-balanced?"

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"Yes, they are group-balanced."

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"Yes, they are group-balanced."

### ACADEMIC

"Your proposal is not optimal for you by your own preferences, regardless of how you tradeoff fairness and accuracy."

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"If you care sufficiently about fairness relative to accuracy, then your proposal **may be optimal** for your goals."

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why might X be group-balanced?

- suppose X has a group-dependent meanings
- e.g., frequent moves signal high creditworthiness for high-income borrowers but low creditworthiness for low-income borrowers
- maximizing accuracy for the high-income group leads this group to have the lower error (and vice versa)

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why might X be group-skewed?

- suppose X is asymmetrically informative
- e.g., medical data is recorded more accurately for high-income patients than low-income patients
- best algorithm coincides for both groups and implies a lower error for high-income patients

# generalizations

beyond absolute difference

- results extend when unfairness is measured as  $|\phi(e_r) \phi(e_b)|$ where  $\phi$  is some continuous strictly increasing function
- if  $\phi$  is log, then this corresponds to evaluating fairness using the ratio of errors rather than their difference

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different loss functions for evaluating fairness and accuracy

- qualitative result extends whenever the two loss functions aren't "directly opposed"
- group-balance generalizes to whether  $F_X$  belongs to usual Pareto frontier
  - when this condition is satisfied, then the fairness-accuracy frontier is identical to the usual Pareto frontier
  - otherwise, the fairness-accuracy frontier is the union of the Pareto frontier and a positively-sloped sequence of lines

### special case: when group identity is an input

#### Proposition

if G is an input in X, then  $\mathcal{E}(X)$  is a rectangle with sides parallel to axes and  $\mathcal{F}(X)$  is the line segment from  $R_X = B_X$  to  $F_X$ 



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 frontier is rawlsian: worse-off group gets best feasible error (no matter which optimization problem we solve) what happens when G is added as an input?



## what happens when G is added as an input?



### what happens when G is added as an input?



**result (informal):** for any designer preference (in our permitted class), access to *G* reduces the error for the worse-off group.

• not true for the other group

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- how limiting is this for the designer? are there garblings that can implement the designer's favorite (unconstrained) point?

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- we formulate a problem of "input design":
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  - decision-maker chooses an algorithm (based on this garbling) to maximize accuracy
- how limiting is this for the designer? are there garblings that can implement the designer's favorite (unconstrained) point?
- ask whether the optimal garbling could involve completely banning a covariate (such as group identity)
  - if X is group-balanced, then banning group identity make every designer strictly worse off

huge literature on algorithmic fairness in CS

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- characterize the fairness-accuracy frontier (in contrast to focusing on a specific optimization problem)
- If formulate problem of choosing inputs as one of "information design" (final input-design section)
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recent empirical work in economics:

• Obermeyer, Powers, Vogeli, Mullainathan (2019), Arnold, Dobbie, and Hull (2021), Fuster, Goldsmith-Pinkham, Ramadorai, Walther (2021)

## Conclusion

- framework for evaluating the accuracy/fairness tradeoffs of algorithms
- characterized the fairness-accuracy frontier over different designer preferences for how to trade off these criteria
- explained how certain statistical properties of the algorithm's inputs impact the shape of this frontier
- in some cases (e.g., when the inputs are group-balanced), there are conclusions/policy recommendations that hold **for all** designer preferences in a broad class

thank you

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**result** (informal): banning G uniformly worsens the Pareto frontier if and only if X is group-balanced.



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#### takeaways:

- when X is group-balanced, all designers strictly benefit from allowing the algorithm to condition on G
- this is true even for an Egalitarian designer: disparate treatment may be necessary to remove disparate impact

**result** (informal): banning X' uniformly worsens the Pareto frontier if and only if X' reduces the error for the worse-off group.



(a) X' reduces group b's error

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**takeaway:** active policy debate regarding whether to ban test scores in admissions decisions.

- so long as G is permissible, then excluding test scores makes all designers worse off
- if G is not a permitted input (e.g., California), then it can be strictly optimal to ban X' (see example in paper)
what happens when the designer only controls the inputs?











example: drop an input



#### example: coarsening

Boalt Hall, UC Berkeley's law school, attracted national media attention in 1997 when its entering class of 268 included only one African-American.<sup>17</sup> The following year, administrators made a number of changes to their admissions policy. The new policy no longer assigns candidates Academic Index Scorespreviously a function of undergraduate GPA (weighted by the quality of the candidate's undergraduate institution) and LSAT score. Indeed, it no longer adjusts candidates' GPAs to account for the quality of their undergraduate institutions. Nor does it consider candidates' exact LSAT scores; instead, LSAT scores are partitioned into intervals, and the admissions committee only learns which interval contains the candidate's score.

(Chan and Eyster, 2003)



input design: feasible and pareto sets

let  $f_T$  denote the utilitarian-optimal algorithm given T

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the **feasible set** under **input design** given X is

 $\mathcal{E}^{*}(X) := \{ e(f_{T}) : T \text{ is a garbling of } X \}$ 

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Definition the **feasible set** under **input design** given X is

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the fairness-accuracy frontier under input design given X is

$$\mathcal{F}^{*}(X) := \left\{ e \in \mathcal{E}^{*}(X) : \underbrace{\text{no } e' \in \mathcal{E}^{*}(X) \text{ s.t. } e' \succ_{FA} e}_{e \text{ is FA-undominated in } \mathcal{E}^{*}(X)} \right\}$$

#### how powerful is input design?

• let  $e_0$  be the utilitarian's best payoff given no information



### how powerful is input design?

- let e<sub>0</sub> be the utilitarian's best payoff given no information
- define  $H := \{(e_r, e_b) : p_r e_r + p_b e_b \le e_0\}$



# how powerful is input design?

- let  $e_0$  be the utilitarian's best payoff given no information
- define  $H := \{(e_r, e_b) : p_r e_r + p_b e_b \le e_0\}$
- lemma:  $\mathcal{E}^*(X) = \mathcal{E}(X) \cap H$

(see also Alonso and Camâra, 2016)

 $\implies$  any point that is feasible given X and in the halfspace H can be implemented using some garbling of X



# how powerful are informational constraints?

#### Proposition

(a) If X is group-balanced, then  $\mathcal{F}(X) = \mathcal{F}^{*}(X)$  iff  $R_{X}, B_{X} \in H$ 

(b) If X is g-skewed, then  $\mathcal{F}(X) = \mathcal{F}^{*}(X)$  iff  $G_{X}, F_{X} \in H$ 



**takeaway:** under weak conditions, designer can implement favorite (unconstrained) outcome by designing the algorithmic inputs

# will the designer want to exclude inputs?

# add/ban covariates?

- regulatory question: should certain inputs be banned?
  - some group identities are already banned (e.g. race, religion for hiring or bank loans)
  - other covariates increasingly prohibited due to fairness concerns (e.g. universities excluding test scores)

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- we can study this using our framework
  - X is the permitted part, X' is the input in question
- question: how does the input-design frontier \$\mathcal{F}^\*(X)\$ compare to \$\mathcal{F}^\*(X, X')\$?

- $Y \in \{0,1\}$  with  $P(Y = 1 \mid G = g) = 1/2$  for both groups g
- X is a null signal, while  $X' \in \{0,1\}$  is a binary signal where
  - X' = Y with probability 1 if G = r
  - X' = Y with probability 0.6 if G = b

so X is more informative about type for group r

• the designer is Egalitarian (payoff is  $-|e_r - e_b|$ )

• sending the null signal X leads to a payoff of |0.5 - 0.5| = 0.



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- any information provided about X' will be used by the agent to improve accuracy
- but this information decreases r's error more than b's error, contributing to a larger gap

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- any information provided about X' will be used by the agent to improve accuracy
- but this information decreases r's error more than b's error, contributing to a larger gap
- the designer's payoffs are strictly negative when any information about X' is provided to the agent

# uniform worsening of the frontier

at the other extreme. . .

#### Definition

say that excluding X' given X uniformly worsens the frontier if every point in  $\mathcal{F}^*(X)$  is FA-dominated by a point in  $\mathcal{F}^*(X, X')$ 

- $\hookrightarrow$  every designer strictly prefers to send information about X'
- $\hookrightarrow$  excluding X' cannot be justified by a fairness-accuracy preference

# excluding group identity

first compare X to (X, G)

**result:** suppose  $R_X, B_X \in H$ . excluding G uniformly worsens the frontier if and only if X is group-balanced



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#### takeaways

- when X is group-balanced, all designers benefit from sending some information about G
- conditioning on *G* means applying an information policy that is asymmetric across groups
- result suggests that **disparate treatment** may be necessary to preclude **disparate impact**
- echos previous findings in the statistical discrimination literature (e.g., Chan and Eyster, 2003)

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$$\{1\} = \operatorname*{argmin}_{d \in \mathcal{D}} \mathbb{E}[\ell(a, Y, g) \mid X = x, X' = x', G = g]$$

while

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this is a weak condition:

• says only that the additional information in X' can change the optimal assignment for some individual in group g

**result:** suppose X reveals G, and let X be g-skewed. excluding X' given X uniformly worsens the frontier  $\iff$  X' is

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(a) X' reduces group b's error

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#### takeaways

there is an active debate regarding whether to ban test scores in admissions decisions

since test scores are likely to be decision-relevant for both groups, our result suggests that:

- so long as G is permissible, then excluding test scores makes all designers worse off
- if G is not a permitted input (as is the case in California), then it can be strictly optimal to ban X' (as in previous example)

# nuances/qualifications

- our results depend critically on our assumption that the designer has access to a **fully flexible** garbling of the inputs X
- do not imply a ranking between sending X' (un-garbled) versus excluding it